# A Local Scalable Distributed Expectation Maximization Algorithm for Large Peer-to-Peer Networks

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# Roadmap

### Introduction

### 2 Motivation

3 Problem statement, contribution

### 4 Locality

### 5 Background

- Expectation maximization
- Notations

### 6 P2P EM algorithm

Experimental results

## Conclusion

## Introduction

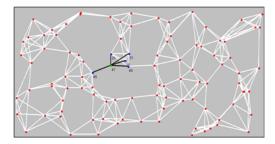


Figure: A P2P network

- Highlights:
  - Highly scalable
  - Asynchronous
  - Completely decentralized
  - Ad-hoc connections

## Introduction

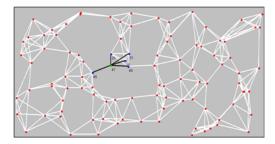


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### Data mining in P2P networks?

- Millions of peers (Skype  $\sim$  50 million)
- Dynamic topology and data peers can join/leave at any time
- No global clock completely asynchronous
- Same features across all peers
- Communication reliable, bandwidth-limited, asynchronous, asymmetric
- Impracticalities / impossibilities
  - global communication
  - global synchronization

# An example



#### Figure: Centralized vs. in-network computation

- EM very useful for variety of data mining tasks
- Can be deployed in P2P networks for
  - clustering
  - anomaly detection
  - target tracking
  - inferencing
- Centralizing data expensive/impractical; collaborative computing *e.g.* cloud computing can harness power of multiple processors/storage

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Can we develop an EM algorithm for P2P networks?

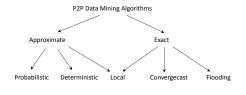


Figure: Distributed data mining algorithms

- S. Datta, K. Bhaduri, C. Giannella, R. Wolff, H. Kargupta. Distributed Data Mining in Peer-to-Peer Networks. IEEE Internet Computing Vol. 10(4), 2006.
- S. Datta, H. Kargupta. Uniform Data Sampling from a Peer-to-Peer Network. ICDCS 2007.
- S. Mukherjee, H. Kargupta. Distributed Probabilistic Inferencing in Sensor Networks using Variational Approximation. JPDC Vol. 68(1), 2008.
- K. Bhaduri, R. Wolff, C. Giannella, H. Kargupta. Distributed Decision Tree Induction in Peer-to-Peer Systems. Statistical Analysis and Data Mining. Vol. 1(2) 2008.

### • Consider large P2P network

- each node has local data which change over time
- each node can exchange messages with immediate neighbors

#### Goal

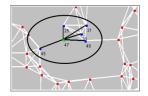
Fit and monitor a gaussian mixture model (gmm) via EM to global data

- Constraints:
  - communication-efficient and scalable
  - asynchronous
  - able to handle dynamic data and network
  - provably correct result compared to centralized computation

- Algorithm for monitoring gmm parameters using EM in large P2P networks
  - local and highly scalable
  - asynchronous
  - provably correct
  - seamlessly handles changes in the data and network

# What is locality?

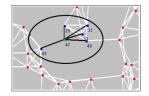
- Every node communicates with only fixed number of other nodes
- Bounded total query size
- Advantages:
  - Scalable
  - Fault-tolerant
  - Robust



#### Figure: Locality of distributed algorithms

# What is locality?

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#### Figure: Locality of distributed algorithms

### Local algorithms

For data dependent algorithms, there exist problem instances whose resource consumption is constant, independent of network size



























 $\pi_1, \overrightarrow{\mu_1}, \mathbf{C}_1$ 

P2P EM



#### Figure: Expectation Maximization



#### Figure: Expectation Maximization

• Given 
$$\mathbf{X} = \{\overrightarrow{x_1}, \overrightarrow{x_2}, \dots, \overrightarrow{x_n}\}$$
, where  $\overrightarrow{x_i} = \mathcal{N}(\overrightarrow{\mu}, \mathbf{C}_s)$ 

- Goal: estimate parameters  $\Theta = \{ \overrightarrow{\mu_1}, \dots, \overrightarrow{\mu_k}, \mathbf{C}_1, \dots, \mathbf{C}_k, \pi_1, \dots, \pi_k \}$
- Approach: maximize log-likelihood of parameters given X

#### Update equations

E-step (estimate the contribution of each point towards each gaussian):

$$q_{s,a} = \frac{\pi_s \mathcal{N}(\overrightarrow{x_a}; \overrightarrow{\mu_s}, \mathbf{C}_s)}{\sum_{r=1}^k \pi_r \mathcal{N}(\overrightarrow{x_a}; \overrightarrow{\mu_r}, \mathbf{C}_r)}$$

no communication

M-step (recompute the parameters of each gaussian):

$$\pi_{s} = \frac{\sum_{a=1}^{n} q_{s,a}}{n}$$

$$\overrightarrow{\mu_{s}} = \frac{\sum_{a=1}^{n} q_{s,a} \overrightarrow{x_{a}}}{\sum_{a=1}^{n} q_{s,a}} \quad \text{communication}$$

$$\mathbf{C}_{s} = \frac{\sum_{a=1}^{n} q_{s,a}(\overrightarrow{x_{a}} - \overrightarrow{\mu}_{s})(\overrightarrow{x_{a}} - \overrightarrow{\mu}_{s})^{\mathrm{T}}}{\sum_{a=1}^{n} q_{s,a}}$$

- $P_1, \ldots, P_p$  a set of peers
- Data stream at P<sub>i</sub>

$$S_i = \left[\overrightarrow{x_{i,1}}, \overrightarrow{x_{i,2}}, \dots, \overrightarrow{x_{i,m_i}}\right]$$

• Global input 
$$\mathcal{G} = \bigcup_{i=1,...,p} S_i$$

•  $X_{i,j}$ : messages sent by  $P_i$  to  $P_j$ 

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### Goal

Build and monitor gmm model on  ${\mathcal G}$  without collecting  ${\mathcal G}$ 

# Thresholding problem

- Problem 1: Compute gmm parameters
  - O(n) communication for exact computation imes
- Problem 2: Given pre-computed parameters, monitoring them vs.  ${\cal G}$ 
  - Less than O(n) communication...very efficient ✓

# Thresholding problem

- Problem 1: Compute gmm parameters
  - O(n) communication for exact computation ×
- Problem 2: Given pre-computed parameters, monitoring them vs.  ${\cal G}$ 
  - Less than O(n) communication...very efficient ✓
- Sufficient statistics
  - Knowledge:  $\mathcal{K}_i = S_i \bigcup_{P_j \in \Gamma_i} X_{j,i}$
  - Agreement:  $A_{i,j} = X_{i,j} \cup X_{j,i}$
  - Withheld:  $\mathcal{W}_{i,j} = \mathcal{K}_i \setminus \mathcal{A}_{i,j}$

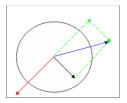


Figure: Set statistics

Conflicting objectives:

- For correct computation,  $\mathcal{K}_i = \mathcal{G}$
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### Solution

- Decompose domain into several non-overlapping convex regions such that any function computed on  ${\cal G}$  remains invariant inside each convex region
- Even if  $\mathcal{K}_i \neq \mathcal{G}$ ,  $\mathcal{F}(\mathcal{K}_i) = \mathcal{F}(\mathcal{G})$  inside any such region
  - Example: Is  $||\mathcal{G}|| < \epsilon$ ?
- Still nobody knows  $\mathcal{G}$ ...

# Local criterion

- Need conditions on local set statistics to infer about  $\ensuremath{\mathcal{G}}$ 

#### Theorem

For each peer and each of its neighbors, if all its set statistics  $\mathcal{K}_i$ ,  $\mathcal{A}_{i,j}$ ,  $\mathcal{W}_{i,j}$  are in same convex region, then so is  $\mathcal{G}$ 

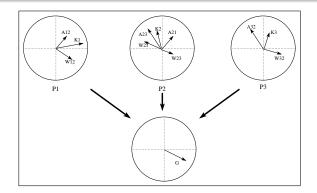


Figure: An example

- Allows a peer to terminate computation and communication whenever stopping condition is satisfied irrespective of other conditions
- Still guarantees eventual correctness
- Remarkably efficient in pruning messages
- Allows a peer to sit idle until an event occurs:
  - send or receive message
  - change in local data
  - change in immediate neighborhood

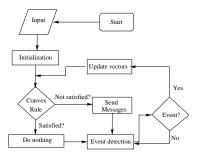
# Back to EM

Monitoring algorithm:

- 1 Input: local dataset, precomputed parameters, error threshold  $\epsilon$
- **2** Goal: monitor  $\mathcal{L}(\Theta)$ ,  $\pi$ ,  $\overrightarrow{\mu}$ , **C**
- 8 Initialization

• 
$$S_i = \left\{ q_{i,s,a} \left( \overrightarrow{x_{i,a}} - \overrightarrow{\overrightarrow{\mu}_s} \right) \right\}$$

- Compute sufficient statistics vectors
- Define convex regions



#### Figure: Flowchart of algorithm

# Computing EM models

- Monitoring algorithm raises an alarm on correct detection
- For closed-loop solution, sample data, rebuild model
- Non-local solution correctness of monitoring algorithm minimizes false dismissals and false alarms



- Simulated data consists of multivariate correlated gaussians with arbitrary parameters
- · Parameters changed at fixed simulator intervals

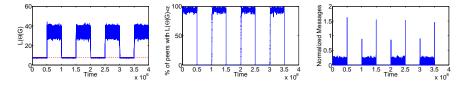


Figure: Experimental results in monitoring mode

# Closed loop results

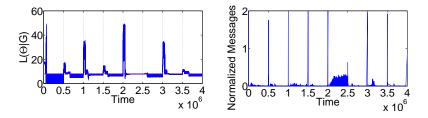


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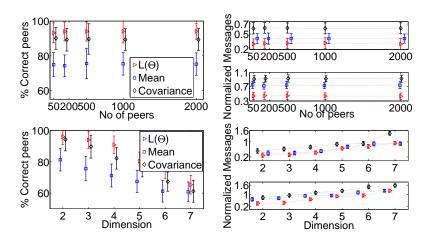


Figure: Scalability results

- First work on developing a local algorithm for gmm monitoring
- Algorithm provably correct, communication efficient, highly scalable, in-network and asynchronous
- Extensive experimental results show low communication cost and correctness of results

#### **Resources:**

- http://ti.arc.nasa.gov/profile/kbhaduri/
- Distributed Data Mining Bibliography: http://www.csee.umbc.edu/~hillol/DDMBIB/